CSCM12 – Software concepts and efficiency   March 10th 2025
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Lab 5: stacks, priority queues and heaps

This week, you should try to get to do question 1. Question 2 is a **challenge** task that might be interesting to revisit later on, after the lecture next week on trees, but it should be accessible today; the first half should be straightforward and not too different in spirit from question 1.

1. **Implementing stacks** Open the file `StackExercises.java` available on canvas. You can see we have defined there a couple of interfaces for stacks:

   - A minimal one called `MyStackMinimal<T>` that says that a stack should contain the push and pop operations

   - A slightly more expanded one called `MyStack<T>` with a couple extra operations.

   (a) Extend `MyArrayStackInt` to an implementation of `MyStack<Integer>`. Use that to run the commented code in the main function (if you want to know what that does, this is the algorithm discussed in the lecture slides that I skipped over).

   (b) Taking inspiration from `MyArrayStackInt`, provide another implementation of `MyStack<T>` based on a linked list structure. You should use the `LinkedList<T>` class of java[1].

   Test that your implementation works by replacing the array-based stack in the main function with your own class.

2. **Using stacks** Now, we will see how to use stacks for another algorithmic problem. Inputs will be strings that can contain

   - Either letters or spaces
   - Or parentheses of several kind:

   $$\texttt{'(' ')' '\{' '\}' '[' ']' '<' '>'}$$

   Write a method

   ```
   static boolean isWellParenthesized(String str)
   ```

   that says whether a string is well-parenthesized or not. The goal will be to recognize strings that are well-parenthesized in the obvious way. Here are a couple of examples of well-parenthesized strings:

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[1]C.f. https://docs.oracle.com/en/java/javase/19/docs/api/java.base/java/util/LinkedList.html

```
"aa( aaa aa aa) { ( adcas ) x (af ( fa )ss )[asf  ]afa }aaa"
"<(<<[]> [ [] ][]>)>"
"()()aa(bb)()"
"aaaa"
```

Here are some strings which are not well-parenthesized:

```
"((())"
"<)"
")(())"
```

3. **Priority Queues (Challenge)** Now we are going to look at priority queues. Open the file `PriorityQueuesExercises.java` from canvas and take some time to recall the interface.

   (a) Fill the code of the class `ListPriorityQueue<T>` that aim to implement a list-based implementation of a priority queue (using an attribute of type `LinkedList<T>`). What are the complexity of the operations you have implemented?

   (b) Now the rest of the exercise is concerned with implementing a priority queue using a heap structure.

   Heaps are structure that are easier to picture using tree-like strutures, but which actually happen to be well-suited to array representations, because the underlying trees are balanced in a very rigid sense.

   So mathematically, a heap will be a kind of labelled tree, where the labels contain an integer priority. We furthermore impose two constraints on those structures:

   - The priority of any node should be inferior or equal to the priority of its children
   - Below any node in the structure, the depth of the subheap on the left-hand side should be the same as the one on the right-hand side, or one level deeper.

   In Figure 1, you have some example of heaps and non-heaps, where only the priorities are used to label the nodes.

   Because of the heap structure, we know that the root contains an element with highest priority, so dequeuing will correspond to deleting the root. Doing so however breaks the heap structure; thankfully, there is a way to repair that using only $\mathcal{O}(\log(n))$ number of steps (where n is the number of nodes in a heap).

   Similarly, to enqueue an element, we shall simply add a node at the most natural spot, and then repair the heap structure in $\mathcal{O}(\log(n))$ steps.

   These two operations, enqueuing and dequeueing, are depicted in Figures 2 and 3 respectively. Enqueuing should be consulted first, as the
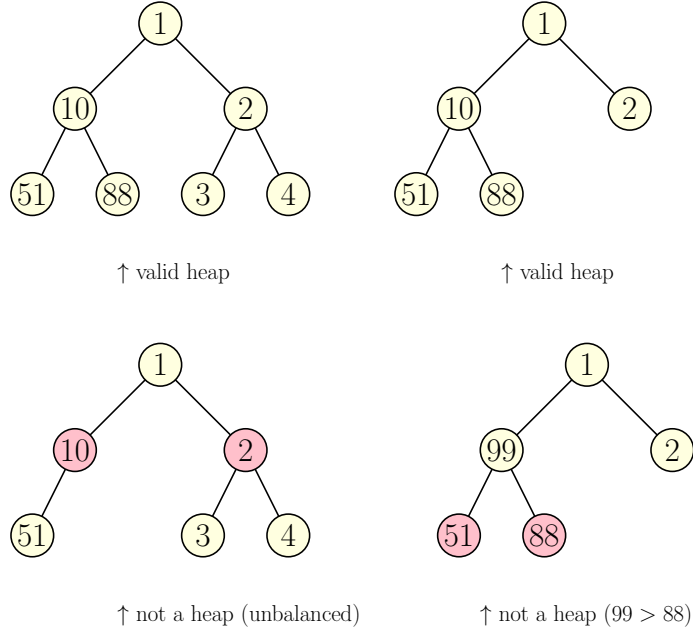
Figure 1: Some example of heaps and non-heaps

repair procedure for dequeueing acually calls the procedure for enqueuing.

Now in the provided code, the heaps are represented in dynamic arrays. The idea is that the label of every node is contained in a cell, and that the tree structure is implicit in the indexing:

- The root corresponds to the cell at index 0
- Regarding the node represented by the cell at index i, its left child is at index $2i + 1$ and its right child at cell $2i + 2$ (if they exist)
- The parent of the node represented by the cell at index j is at index $\left\lfloor \frac{j-1}{2} \right\rfloor$

With these notations, the two heaps in Figure 1 are represented by the following arrays:

```
{ 1, 10, 2, 51, 88, 3, 4 }
{ 1, 10, 2, 51, 88}
```

And all of the tree-like structures in Figure 2 may be represented by

```
{ 1, 10, 8, 51, 88, 11, 9, 99, 88 }
{ 1, 10, 8, 51, 88, 11, 9, 99, 88, 7 }
{ 1, 10, 8, 51, 7, 11, 9, 99, 88, 88 }
{ 1, 7, 8, 51, 10, 11, 9, 99, 88, 88 }
```

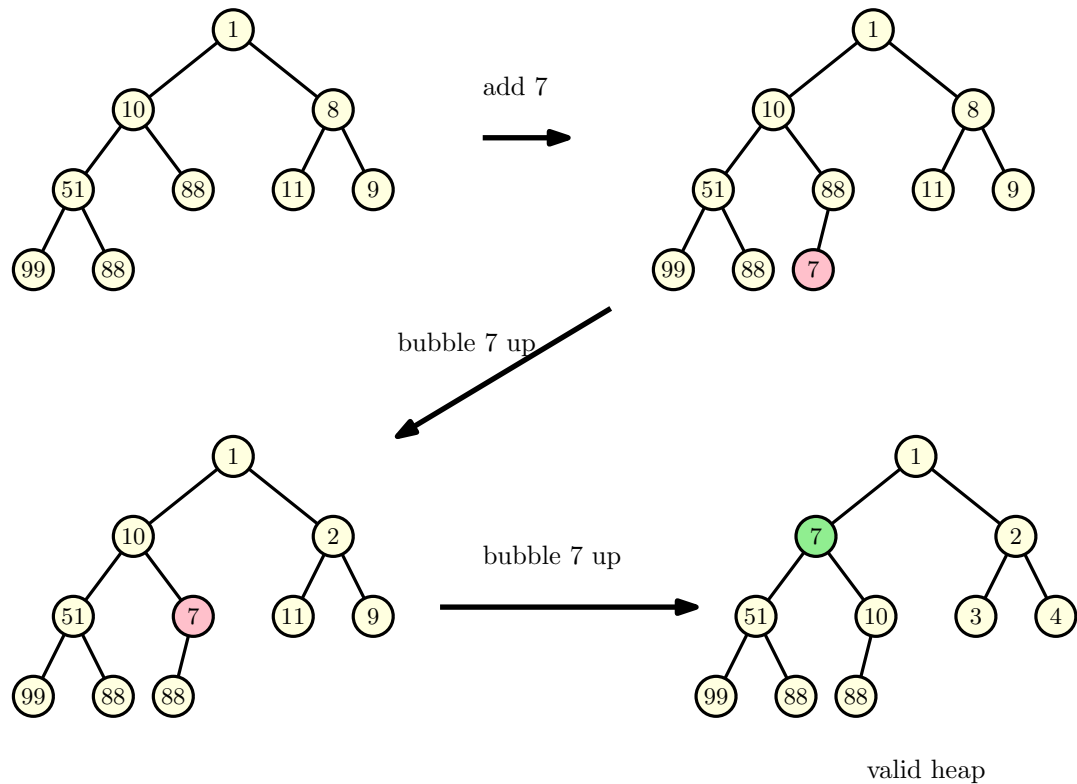Now the code for the heap-based queue is already somewhat pre-filled in

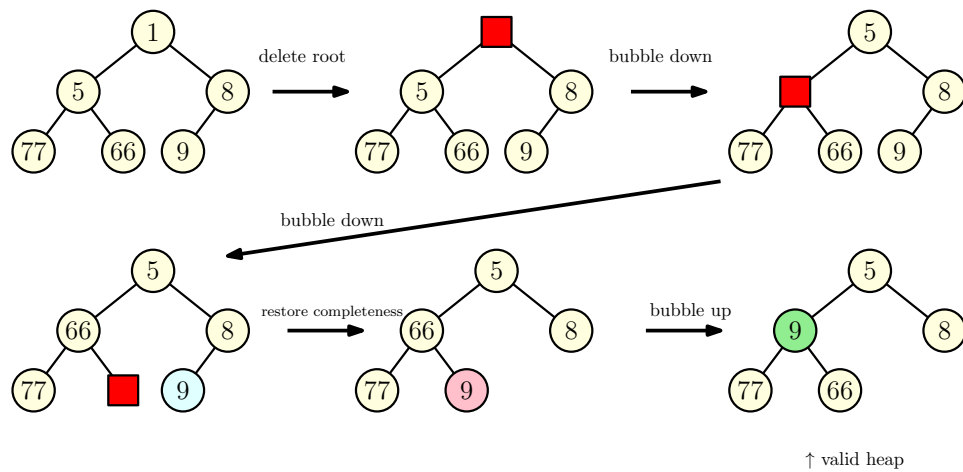Figure 2: Inserting a new element with priority 9 in a heap and repairing.

Figure 3: Removing the root of a heap and repairing. Note that the procedure is in two steps: first a "hole" is propagated down, and then, in order to restore the second invariant, the deepest rightmost node is put in place of the newly created hole, and then bubbled up using the same repair procedure as for enqueuing.

the file, except for the part that handle the repairs, handled by private methods

```
private void bubbleUp(int index);
private int  bubbleDown(int index);
```

Your task is to complete the code of these functions to obtain a working dequeue implementation as sketched in the handout!