

CS-205 lecture 11:

Interactive programming and further topics

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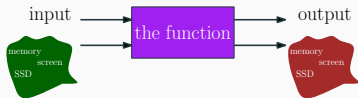
Autograder is live

Link to submit your coursework

<https://csautograder.swansea.ac.uk/web/project/69>

- Detailed submission instructions on canvas
- A bug was reported to me yesterday, should be fixed
- All tests re-ran after final submission
- 37 marks allocated automatically, the rest by handgrading
- The handgrader might compensate for harsh automated grading

Last week: interactive programming



What we have seen

How to

- write types for programs with IO side-effects in types **IO** a
`print :: Show a => a -> IO ()`
`getLine :: IO String`
- combine them using bind `>>=` or the **do** notation
- a couple of examples
- compile haskell programs using `ghc`

Extended example: validating input

```
getYN :: String -> IO Bool
getYN prompt = putStr (prompt ++ "[y/N]:") >>
  hFlush stdout >>
  getLine >>= \s ->
  if s `elem` answers then
    return (s `elem` yanswers)
  else
    putStrLn "Wrong input!" >> getYN prompt
where yanswers = ["y", "Y", "yes", "Yes"]
      answers  = yanswers ++ ["", "n", "N", "no", "No"]
```

Extended example: validating input

```
getYN :: String -> IO Bool
getYN prompt = do {
    putStr (prompt ++ "[y/N]:");
    hFlush stdout;
    s <- getLine;
    if s `elem` answers then
        return (s `elem` yanswers)
    else
        do {
            putStrLn "Wrong input!";
            getYN prompt
        }
}
where yanswers = ["y", "Y", "yes", "Yes"]
      answers  = yanswers ++ ["", "n", "N", "no", "No"]
```

Some tps/considerations for the lab

- I have not gone over all the **IO** primitives
→ use the online documentation (hackage/hoogle)
- You might need some import statement to import functions like `hFlush` or `isDigit` as in e.g.

```
import Data.Char (isDigit) -- imports only isDigit
import System.IO  -- imports everything in the module
```
- `hFlush stdout` \equiv `fflush(stdout)`
flushes the stdout buffer → forces printing

Some further topics

Warning

The rest of the lecture will survey some topics you could look into if you want to keep writing Haskell in the future/are curious

Before we move on, questions about Haskell/CW/etc?

Ofc you are free to ask at any later point :)

(more detailed explanation on the material below in lecture11.hs)

Further topic 1: monads

Extended do notation?

```
doList :: [(Int,Char)]
doList = do {
    x <- [1..5];
    y <- ['a', 'z'];
    return (x,y)
}
```

Extended do notation?

```
divMaybe :: Int -> Int -> Maybe Int
```

```
divMaybe x 0 = Nothing
```

```
divMaybe x y = Just (x `div` y)
```

```
doMaybe :: Int -> Int -> Int -> Maybe Int
```

```
doMaybe x y z = do {
```

```
    a <- divMaybe x y;
```

```
    b <- divMaybe z a;
```

```
    return (a + b)
```

```
}
```

Workhorse behind this: the Monad typeclass

```
(>>=) :: Monad m => m a -> (a -> m b) -> m b  
return :: Monad m => a -> m a
```

- $m :: * \rightarrow *$ is a variable, but not for a type
- Monads consist of a very generic, yet useful abstractions
- Typical instances: **Monad IO**, **Monad []**, **Monad Maybe**,
Monad (Cont r), **Monad (State s)**
- **Monad (State s)** = code “as if” we had mutable variables

Maybe the next example to look at if you are interested

The state monad

```
data State s a = Stateful (s -> a * s)
```

```
return :: a -> State s a
```

```
return x m = (x, m)
```

```
(>>=) :: State s a -> (a -> State s b) -> State s b
```

```
(Stateful c) >>= f = Stateful
```

```
  \m -> let (x, m') = c m in
```

```
    let Stateful g = f x in
```

```
      g m'
```

```
escape :: State s a -> s -> a
```

```
escape (Stateful c) m = fst (c m)
```

- Usefulness wrt IO: one can go back to pure computations via `escape`

(function of type `IO a -> a` named `unsafePerformIO`)

Laziness

Lazy evaluation

With FP languages, there are two popular kind of evaluation strategies

$(\lambda x \rightarrow x + x) (1+2)$

- **Eager/CBV**: evaluate arguments first

$$\begin{aligned}(\lambda x \rightarrow x + x)(1 + 2) &\rightarrow (\lambda x \rightarrow x + x)3 \\ &\rightarrow (\lambda x \rightarrow x + x)3 \\ &\rightarrow 3 + 3 \\ &\rightarrow 6\end{aligned}$$

- **Lazy/CBN**: substitute arguments in the function body first

$$\begin{aligned}(\lambda x \rightarrow x + x)(1 + 2) &\rightarrow (1 + 2) + (1 + 2) \\ &\rightarrow 3 + (1 + 2) \\ &\rightarrow 3 + 3 \\ &\rightarrow 6\end{aligned}$$

In pure functional programming languages, the evaluation strategy mostly does not matter for the result!

- Haskell is **lazy**. (there are pros/cons with that)
- It tries to avoid to duplicate computations (call by need strategy)
(as)

$$\begin{aligned}(\lambda x \rightarrow x + x)(1 + 2) &\rightarrow (1 + 2) + (1 + 2) \\ &\rightarrow 3 + 3 \\ &\rightarrow 6\end{aligned}$$

Pro/Cons laziness

Pros:

- Call-by-need can save some shared computation at low intellectual cost
→ nice for rapid prototyping of complicated code
- Some nice idiosyncratic applications in the next slide

Cons:

- Harder to reason about complexity
- Counter-intuitive
- More complicated runtime because thunking is necessary
- `unsafePerformIO` has really hard-to-predict behaviours
- laziness can easily be emulated in eager languages

(essentially replace `a by () -> a`)

Some applications

- Infinite values can be used seamlessly in the language

```
allNats :: [Int]
```

```
allNats = 0 : map (+1) allNats
```

```
-- >>> take 5 allNats
```

```
-- [0,1,2,3,4]
```

- Nice tricks, like support for memoization/dynamic programming **without side-effects or state monad**

Not possible in eager FP languages

Next slides: explanation of the dynamic programming example in `lecture11.hs`

Extended example: binomial (1/4)

Problem

Compute the number of ways $\binom{n}{k}$ to pick k elements among n .

$$\binom{4}{2} = \#\left\{ \begin{array}{c} \text{●} \\ \text{●} \\ \text{○} \\ \text{○} \end{array}, \begin{array}{c} \text{○} \\ \text{○} \\ \text{●} \\ \text{○} \end{array}, \begin{array}{c} \text{○} \\ \text{○} \\ \text{●} \\ \text{●} \end{array}, \begin{array}{c} \text{○} \\ \text{○} \\ \text{○} \\ \text{●} \end{array}, \begin{array}{c} \text{○} \\ \text{○} \\ \text{○} \\ \text{●} \end{array}, \begin{array}{c} \text{○} \\ \text{○} \\ \text{○} \\ \text{●} \end{array} \right\}$$

Extended example: binomial (1/4)

Problem

Compute the number of ways $\binom{n}{k}$ to pick k elements among n .

$$\binom{4}{2} = \#\left\{ \begin{array}{c} \text{⦿} \\ \text{⦿} \\ \text{⦿} \\ \text{⦿} \end{array} \right\}$$

$$\binom{n}{k} = \#\{X \subseteq \{1, \dots, n\} \mid \#X = k\} = \frac{n!}{k!(n-k)!}$$

Extended example: binomial (1/4)

Problem

Compute the number of ways $\binom{n}{k}$ to pick k elements among n .

$$\binom{4}{2} = \#\left\{ \begin{array}{c} \text{⦿} \\ \text{⦿} \end{array}, \begin{array}{c} \text{⦿} \\ \text{⦿} \end{array}, \begin{array}{c} \text{⦿} \\ \text{⦿} \end{array}, \begin{array}{c} \text{⦿} \\ \text{⦿} \end{array}, \begin{array}{c} \text{⦿} \\ \text{⦿} \end{array}, \begin{array}{c} \text{⦿} \\ \text{⦿} \end{array} \right\}$$

$$\binom{n}{k} = \#\{X \subseteq \{1, \dots, n\} \mid \#X = k\} = \frac{n!}{k!(n-k)!}$$

Issue with the closed formula: $n!$ overflows fast while $\binom{k}{n}$ is polynomial if $k = O(1)$.

Alternative way of computing?

Extended example: binomial (2/4)

Decomposition by fixing an element and asking whether it is picked or not.

$$\begin{aligned} \binom{4}{2} &= \# \left\{ \begin{array}{c} \text{●} \\ \text{●} \\ \text{○} \\ \text{○} \end{array}, \begin{array}{c} \text{●} \\ \text{○} \\ \text{●} \\ \text{○} \end{array}, \begin{array}{c} \text{○} \\ \text{●} \\ \text{○} \\ \text{●} \end{array} \right\} + \# \left\{ \begin{array}{c} \text{○} \\ \text{○} \\ \text{●} \\ \text{●} \end{array}, \begin{array}{c} \text{○} \\ \text{○} \\ \text{●} \\ \text{●} \end{array}, \begin{array}{c} \text{○} \\ \text{○} \\ \text{●} \\ \text{●} \end{array} \right\} \\ &= \# \left\{ \begin{array}{c} \text{○} \\ \text{○} \\ \text{●} \\ \text{●} \end{array}, \begin{array}{c} \text{○} \\ \text{○} \\ \text{●} \\ \text{●} \end{array}, \begin{array}{c} \text{○} \\ \text{○} \\ \text{●} \\ \text{●} \end{array} \right\} + \# \left\{ \begin{array}{c} \text{○} \\ \text{○} \\ \text{●} \\ \text{●} \end{array}, \begin{array}{c} \text{○} \\ \text{○} \\ \text{●} \\ \text{●} \end{array}, \begin{array}{c} \text{○} \\ \text{○} \\ \text{●} \\ \text{●} \end{array} \right\} \\ &= \binom{3}{1} + \binom{3}{2} \end{aligned}$$

Extended example: binomial (2/4)

Decomposition by fixing an element and asking whether it is picked or not.

$$\begin{aligned} \binom{4}{2} &= \# \left\{ \begin{array}{c} \text{blue} \\ \text{red} \\ \text{white} \end{array} \right\} + \# \left\{ \begin{array}{c} \text{red} \\ \text{white} \\ \text{blue} \end{array} \right\} \\ &= \# \left\{ \begin{array}{c} \text{red} \\ \text{white} \\ \text{blue} \end{array} \right\} + \# \left\{ \begin{array}{c} \text{blue} \\ \text{white} \\ \text{red} \end{array} \right\} \\ &= \binom{3}{1} + \binom{3}{2} \end{aligned}$$

$$\binom{n}{k} = \binom{n-1}{k-1} + \binom{n-1}{k}$$

Extended example: binomial (3/4)

```
binom :: Int -> Int -> Int
```

```
binom k n | k > n = 0
```

```
binom 0 n      = 1
```

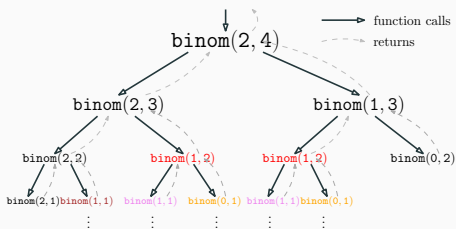
```
binom k n      = binom (k-1) (n-1) + binom k (n-1)
```

Proof of termination: by induction over n .

Extended example: binomial (3/5)

Issue: exponential number of calls

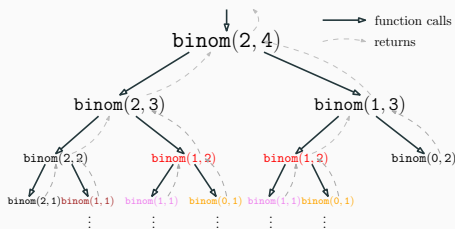
(inefficient)



Extended example: binomial (3/5)

Issue: exponential number of calls

(inefficient)



But there are redundant calls!

- Dynamic programming/memoization: cache the common subcomputations!

Extended example: binomial (4/4)

```
binomial :: Int -> Int -> Int
binomial k n | k > n      = 0
             | otherwise = a ! (k, n)
             where a = array ((0,0),(k,n))
                           [ ((i,j), b i j) | i <- [0..k],
                                               j <- [0..n]]

             b 0 k = 1
             b i j | i == j = 1
             b i j = (a ! (i,j-1)) + (a ! (i-1,j-1))
```

Some caveats:

- The imperative implementation might be more straightforward
- Also does not mesh well with **hash-consing** if the input domain is more complex

Innovation compared to previous years

We are dropping the mandatory verification part from the exam.

Leaves us one extra session. Rough ideas:

- Set up a “real” project with cabal
- Haskell & contemporary FP
 - Other FP languages
 - FP features in traditionally imperative languages
 - Proof assistants
- Some other ideas for further topics you may want to look at on your own that could make use of the module content?
- Q&A, AMA related to the module content
 - (in which case, it would be useful to have questions in advance)
- Am also open to suggestions until Thursday!